

Towards General-Purpose Proof Automation for Lean

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“Towards”



Architecture

Customisation

Debugging

Proof Search

1. Simplify everywhere.
2. Apply safe resolution rules.
3. Apply unsafe resolution rules.

$A B : \text{set } \alpha$

$\vdash A \cap B \subseteq A \cup B$

$A \ B : \text{set } \alpha$

$\vdash A \cap B \subseteq A \cup B$

...

$\vdash (\lambda x, A \ x \wedge B \ x) \subseteq (\lambda x, A \ x \vee B \ x)$

$A \ B : \text{set } \alpha$

$\vdash A \cap B \subseteq A \cup B$

...

$\vdash (\lambda x, A \ x \wedge B \ x) \subseteq (\lambda x, A \ x \vee B \ x)$

...

$\vdash \forall x, A \ x \wedge B \ x \rightarrow A \ x \vee B \ x$

...
 $\vdash \forall x, A x \wedge B x \rightarrow A x \vee B x$

...

$\vdash \forall x, A x \wedge B x \rightarrow A x \vee B x$

...

$x : \alpha$

$h : A x \wedge B x$

$\vdash A x \vee B x$

...

$\vdash \forall x, A x \wedge B x \rightarrow A x \vee B x$

...

$x : \alpha$

$h : A x \wedge B x$

$\vdash A x \vee B x$

...

$h_1 : A x$

$h_2 : B x$

$\vdash A x \vee B x$

...
 $\vdash \forall x, A x \wedge B x \rightarrow A x \vee B x$

...
 $x : \alpha$
 $h : A x \wedge B x$
 $\vdash A x \vee B x$

...
 $h_1 : A x$
 $h_2 : B x$
 $\vdash A x \vee B x$

... | ...
 $\vdash A x$ | $\vdash B x$

```
@[backward unsafe]
lemma or_intro_left  { $\alpha$   $\beta$ } :  $\alpha \rightarrow \alpha \vee \beta$ 
```

```
@[backward unsafe]
lemma or_intro_right { $\alpha$   $\beta$ } :  $\beta \rightarrow \alpha \vee \beta$ 
```

Customisation

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```
@[backward safe]
meta def continuous_comp_rule : tactic unit :=
/- Apply continuous.comp unless the function
   is constant or the identity. -/
```

Customisation

```
@[backward safe]
meta def continuous_comp_rule : tactic unit :=
/- Apply continuous.comp unless the function
   is constant or the identity. -/
```

```
@[backward finisher]
meta def arithmetic_rule : tactic unit :=
/- If target is  $n < m$  or  $n > m$  or ...,
   try linearith and nlinearith. -/
```

Debugging

Debugging

$$\alpha \ \beta \ \gamma : \text{Prop}$$
$$\vdash \alpha \rightarrow (\beta \vee \gamma) \vee \alpha$$

Debugging

$\alpha \beta \gamma : \text{Prop}$
 $\vdash \alpha \rightarrow (\beta \vee \gamma) \vee \alpha$

- intros
- apply or.intro_left
- apply or.intro_right
 - assumption

- intros
- apply or.intro_left
 - apply or.intro_left
 - failure: no rules applicable
 - apply or.intro_right
 - failure: no rules applicable
- apply or.intro_right
 - assumption

- intros
- apply or.intro_left
- apply or.intro_right
- $\alpha \beta \gamma : Prop$
- $a : \alpha$
- $\vdash \alpha$
- assumption

- intros (0 ms)
- 42 rules not applicable (10ms)
- apply or.intro_left (0ms + 5ms)
- apply or.intro_right (0 ms)
 - assumption (0 ms)

Evaluation

Power



Performance



Customisability



Transparency



Please Discuss!

- ▶ Proof search based on simplification and resolution rules.
- ▶ Custom rules for heuristics, decision procedures, ...
- ▶ Interactive search tree for debugging.

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Tagging

- ▶ Speculative tagging

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- ▶ Linting (eg `simp_nf`)

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- ▶ Speculative tagging
- ▶ Linting (eg `simp_nf`)
- ▶ Automated refactoring